



Data Structures and Algorithms (12)

Instructor: Ming Zhang

Textbook Authors: Ming Zhang, Tengjiao Wang and Haiyan Zhao

Higher Education Press, 2008.6 (the "Eleventh Five-Year" national planning textbook)

<https://courses.edx.org/courses/PekingX/04830050x/2T2014/>



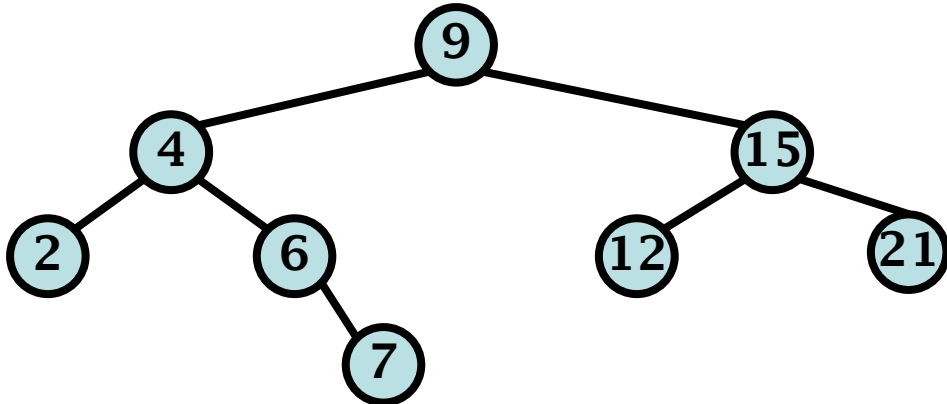
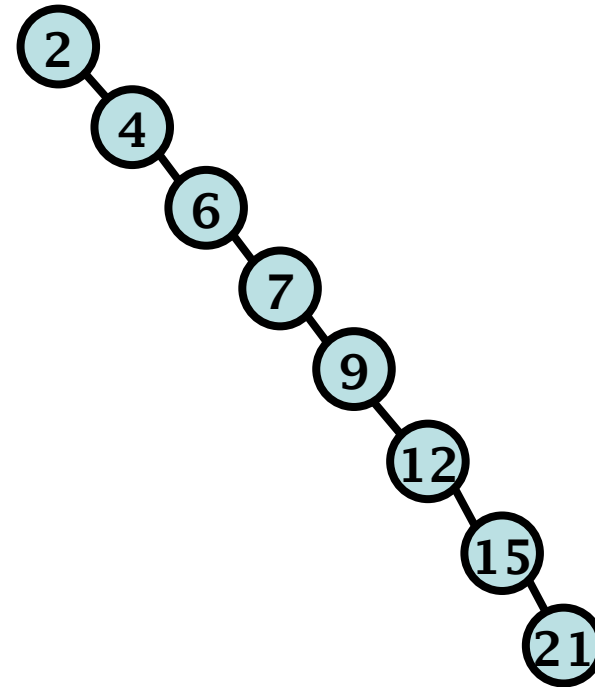
Chapter 12 Advanced data structure

- 12.1 Multidimensional Array
- 12.2 Generalized Lists
- 12.3 Storage management
- 12.4 Trie
- 12.5 Improved binary search tree
 - 12.5.1 Balanced binary search tree
 - Concept and inserting operation of AVL tree
 - Deleting operation and efficiency analysis of AVL tree
 - 12.5.2 Splay Tree

12.5 Improved binary search tree

12.5.1 AVL

- The performance of BST operations are affected by the input sequence
 - Best $O(\log n)$; Worst $O(n)$
- Adelson-Velskii and Landis
 - AVL tree, a balanced binary search tree
 - Always $O(\log n)$

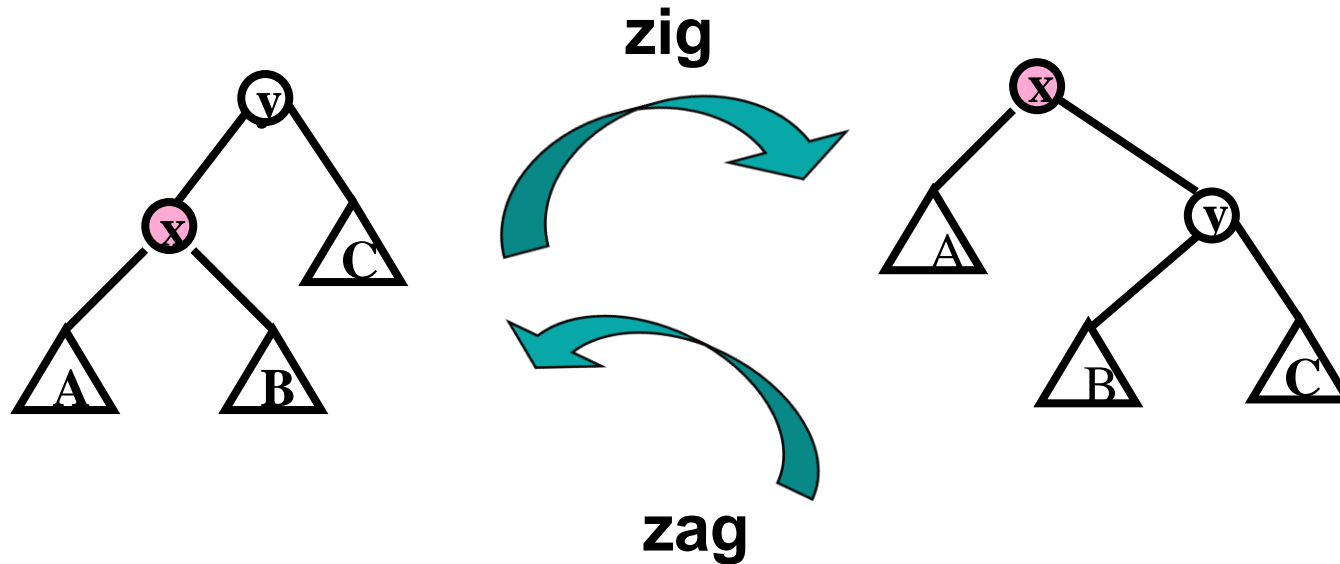


12.5 Improved binary search tree

12.5.1 AVL

• Single Rotation

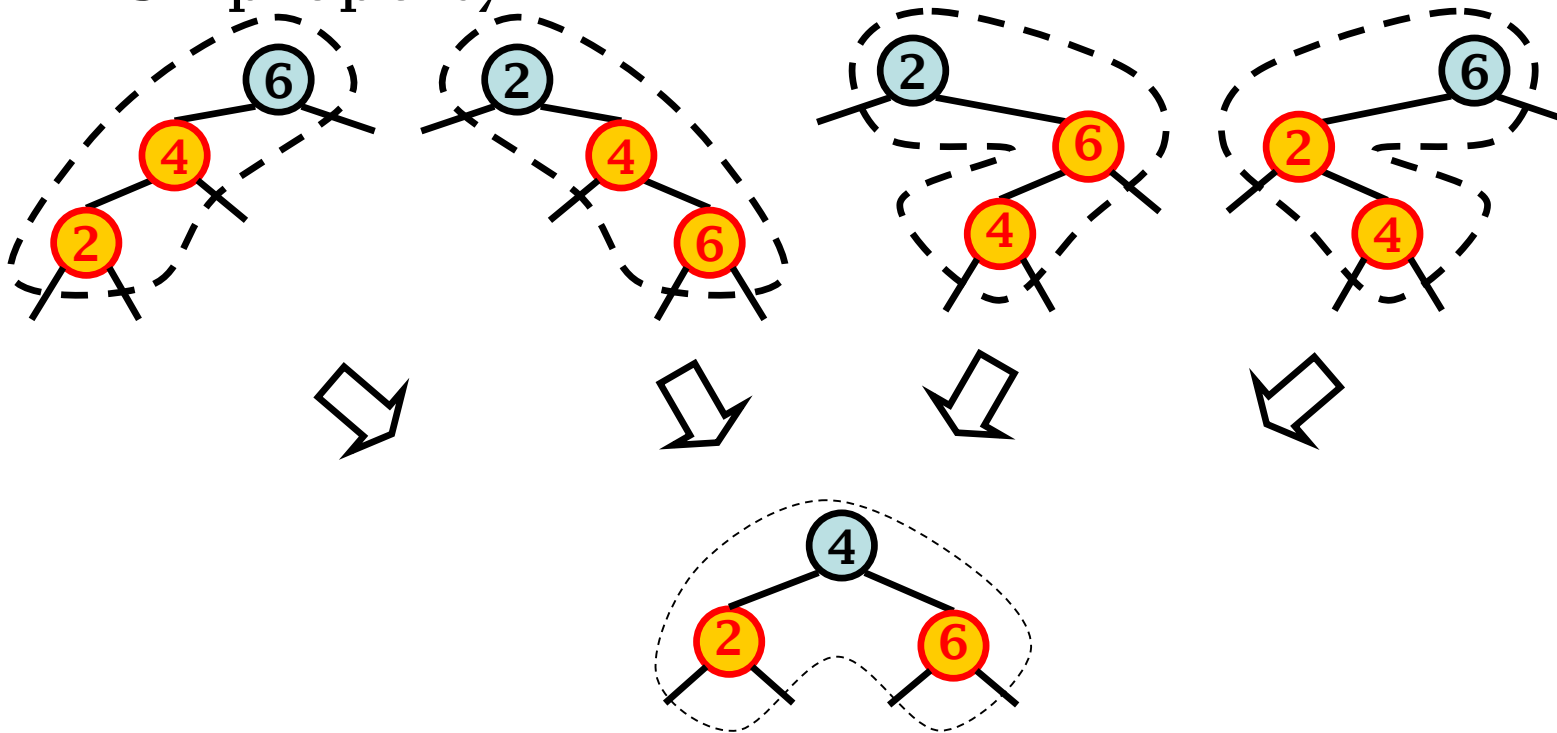
- Swap the node with its father, while keeping the property of BST



12.5 Improved binary search tree

12.5.1 AVL

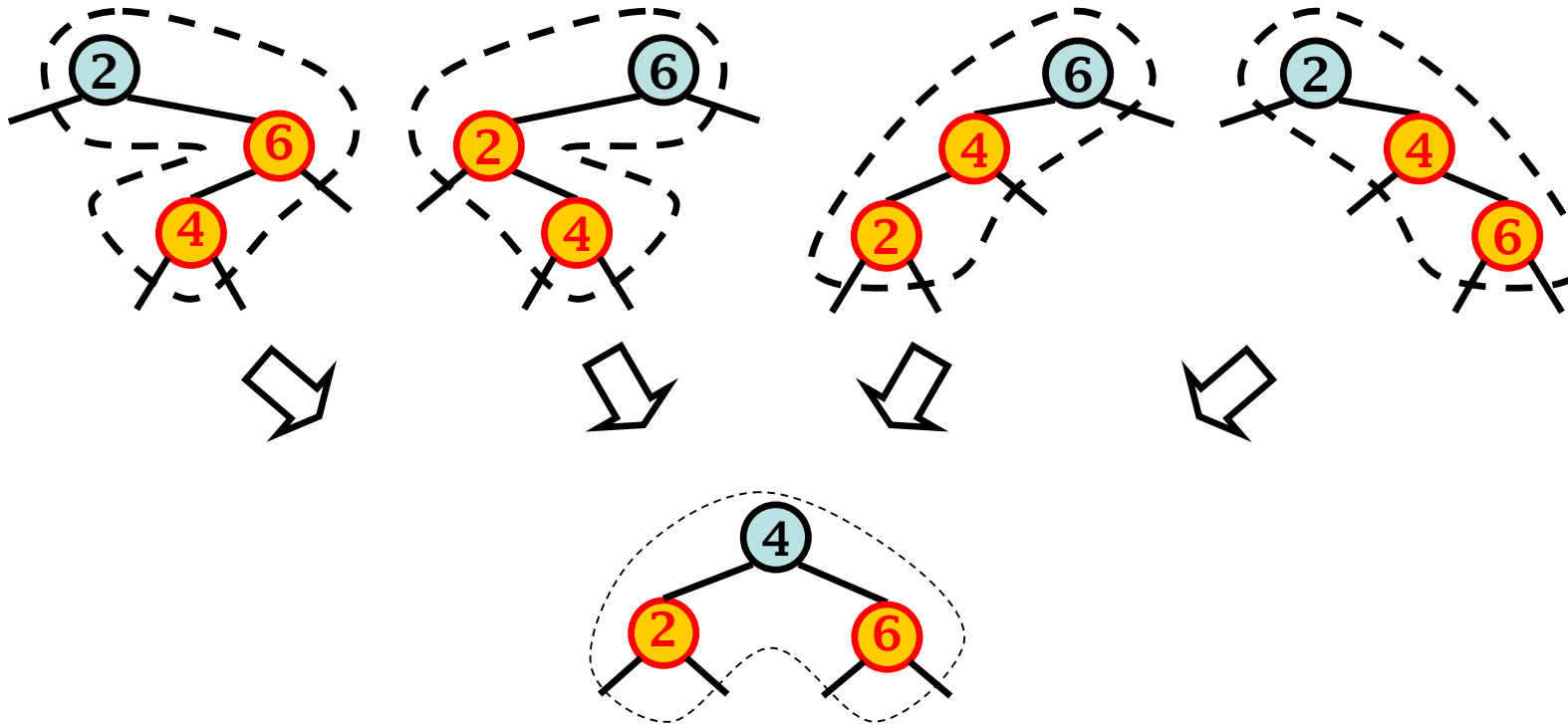
- Single Rotation and Double Rotation: Keep the BST property.



12.5 Improved binary search tree

12.5.1 AVL

- Equivalent rotation: Keep the BST property



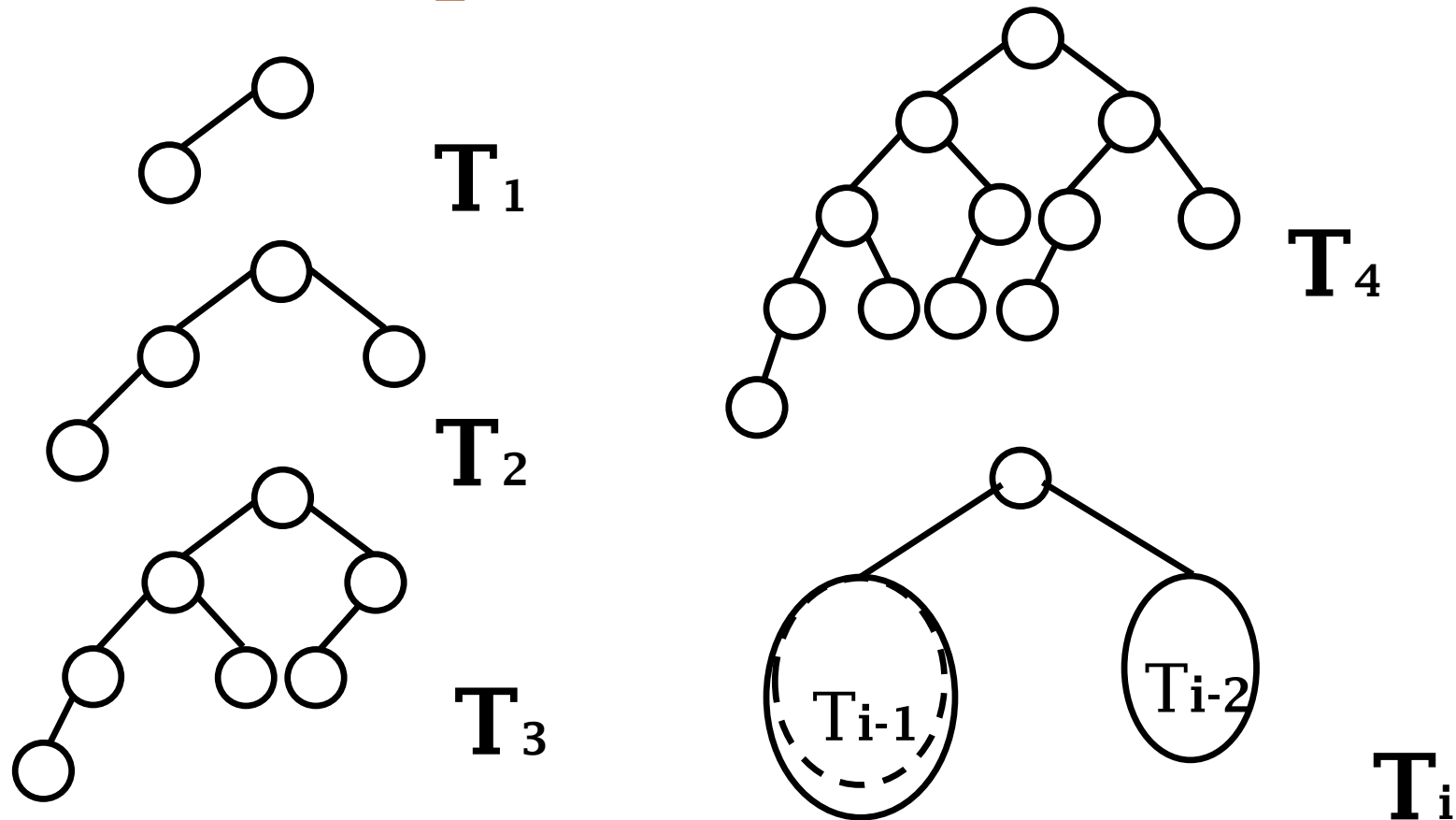
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AVL

- Empty tree is allowed
- The height of AVL tree with n nodes is $O(\log n)$
- If T is an AVL tree
 - Then the left and right subtree of T : T_L , T_R are also AVL trees
 - And $|h_L - h_R| \leq 1$
 - h_L , h_R are the heights of its left and right subtree.

12.5 Improved binary search tree

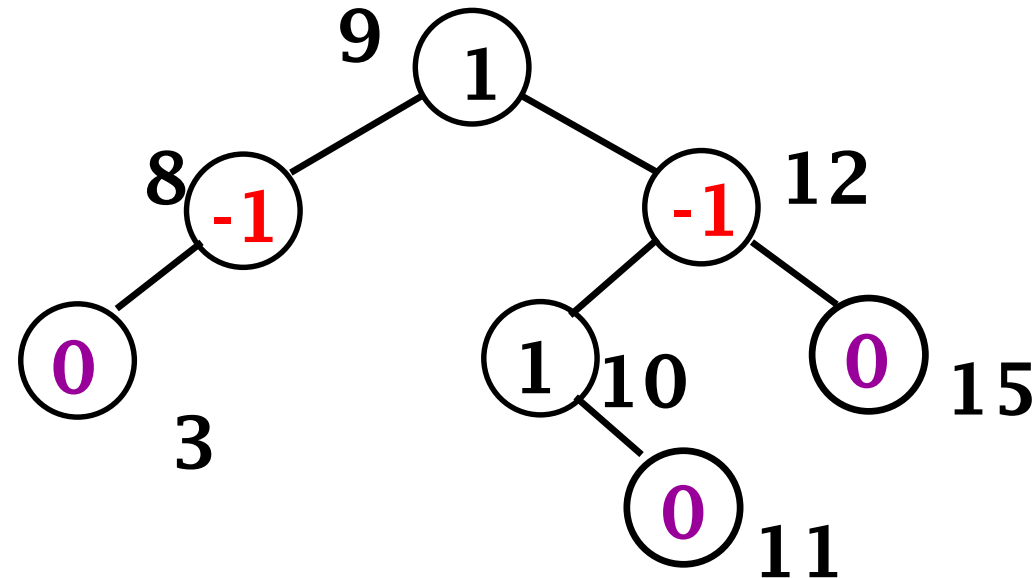
Examples of AVL tree



12.5 Improved binary search tree

Balance Factor

- Balance Factor , $bf(x)$:
 - $bf(x) = height(x_{rchild}) - height(x_{lchild})$
- Balance Factor might be 0, 1 and -1



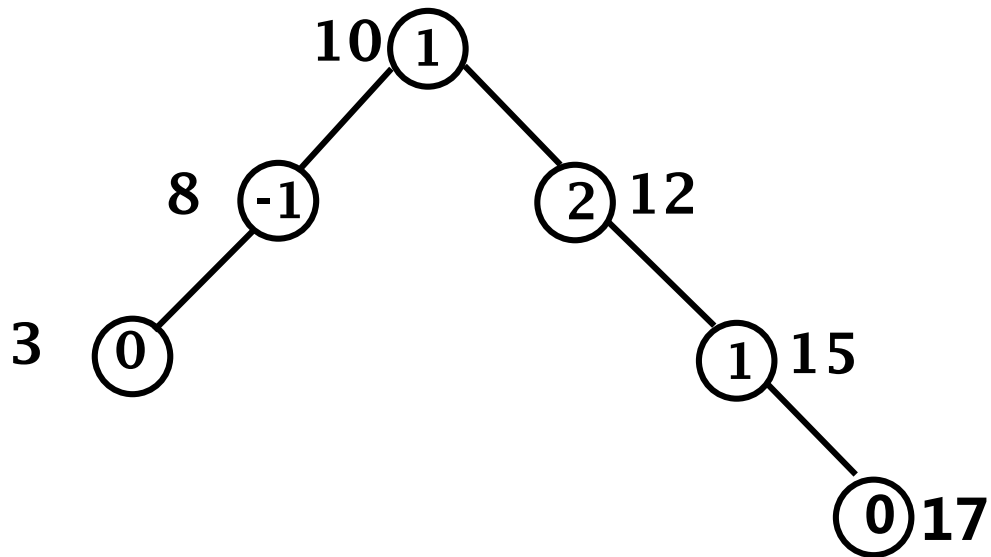
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Insertion in an AVL tree

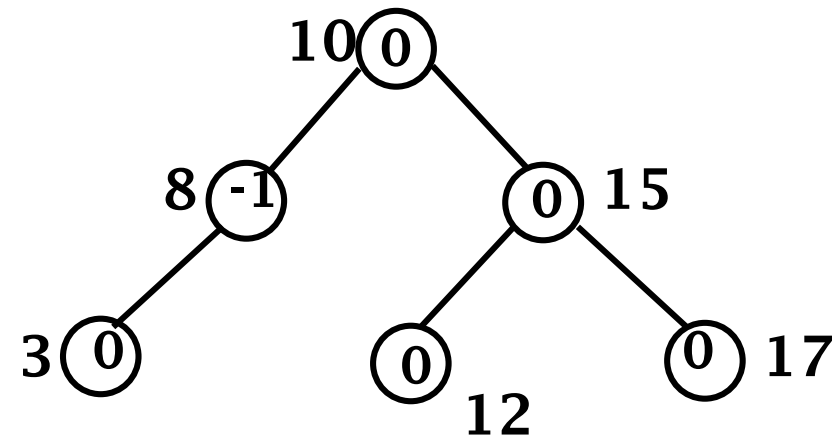
- Just like BST: insert the current node as a leaf node
- Situations during adjustment
 - The current node was balanced. Now its left or right subtree becomes heavier.
 - Modify the balance factor of the current node
 - The current node had a balance factor of ± 1 . Now the current node becomes balanced.
 - Height stays the same. Do not modify.
 - The current node had a balance factor of ± 1 . Now the heavier side becomes heavier
 - Unbalanced
 - “dangerous node”

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Rebalance



Become unbalanced after inserting 17



Adjustment



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- Unbalanced situation occurs after insertion
- Insert the current node as leaf node in BST
- Assume a is the most close node to the current node. And its absolute value of balance factor is not zero.
- The current node s with key is in its left subtree or its right subtree.
- Assume that it is inserted into the right subtree.
The original balance factor:
 - (1) $bf(a) = -1$
 - (2) $bf(a) = 0$
 - (3) $bf(a) = +1$

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- Assume a is the most close node to the current node s . And its absolute value of balance factor is not zero.
 - S is in a 's left subtree or right subtree.
- Assume S is in the right subtree. Because balance factors of nodes in paths from s to a change from 0 to +1. So as for node a :
 1. $bf(a) = -1$, then $bf(a) = 0$, and the height of node a 's subtree stays the same.
 2. $bf(a) = 0$, then $bf(a) = +1$, and the height of node a 's subtree stays the same.
 - Because of the definition of a ($bf(a) \neq 0$), we can know that node a is the root.
 3. $bf(a) = +1$, then $bf(a) = +2$, and adjustment is needed.

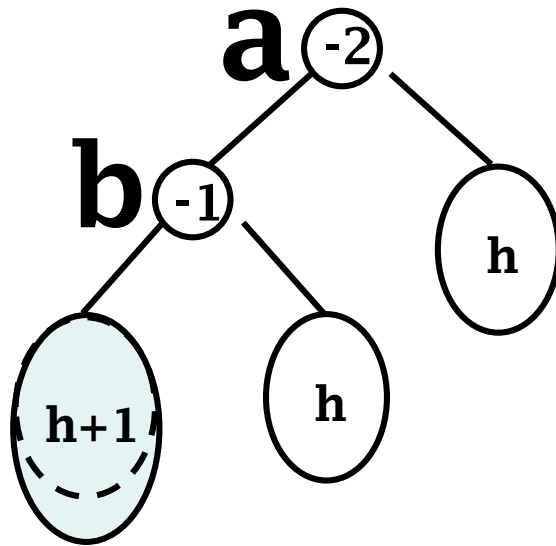
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Unbalanced Cases

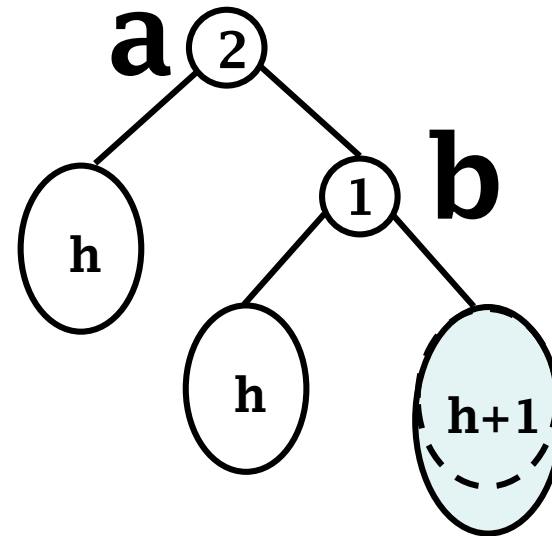
- The balance factors of any nodes must be **0, 1, -1**
- a's left subtree was heavier, $bf(a) = -1$, and $bf(a)$ become -2 after insertion.
 - **LL**: insert into the left subtree of a's left child.
 - Left heavier + left heavier, $bf(a)$ become -2
 - **LR**: insert into the right subtree of a's left child.
 - Left heavier + right heavier, $bf(a)$ become -2
- Likewise, $bf(a) = 1$, and $bf(a)$ become 2 after insertion
 - **RR**: the node that causes unbalanced is in the right subtree of a's right child.
 - **RL**: the node that causes unbalanced is in the left subtree of a's right child.

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Unbalanced Cases



LL



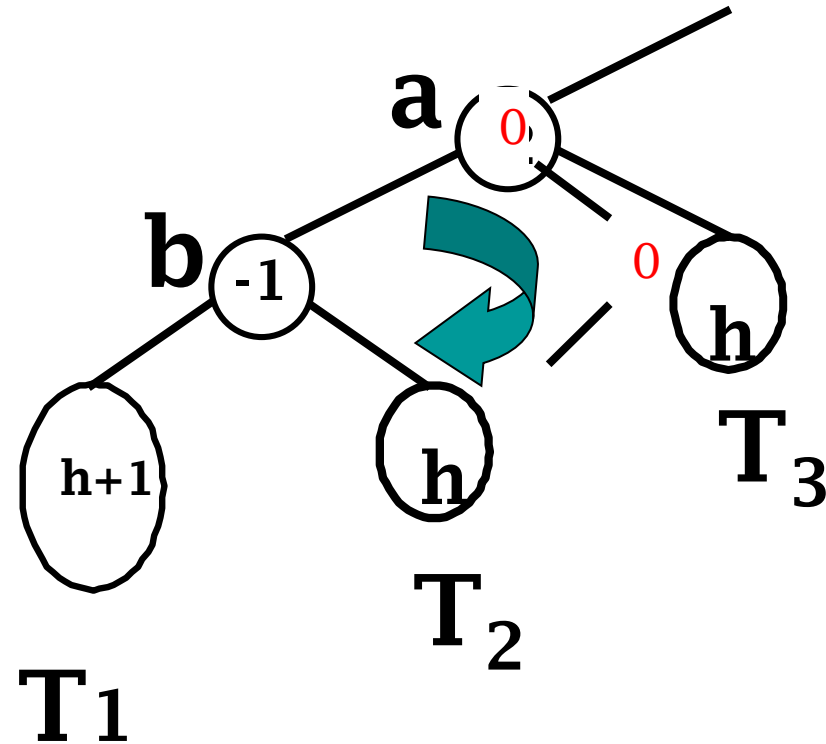
RR

Summary of unbalanced cases

- LL is symmetry with RR, and LR is symmetry with RL.
- Unbalanced nodes happen on the path from inserted node to the root.
- Its balance factor must be 2 or -2
 - If 2, the balance factor before insertion is 1
 - If -2, the balance factor before insertion is -1

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LL single rotation





Insight of Rotations

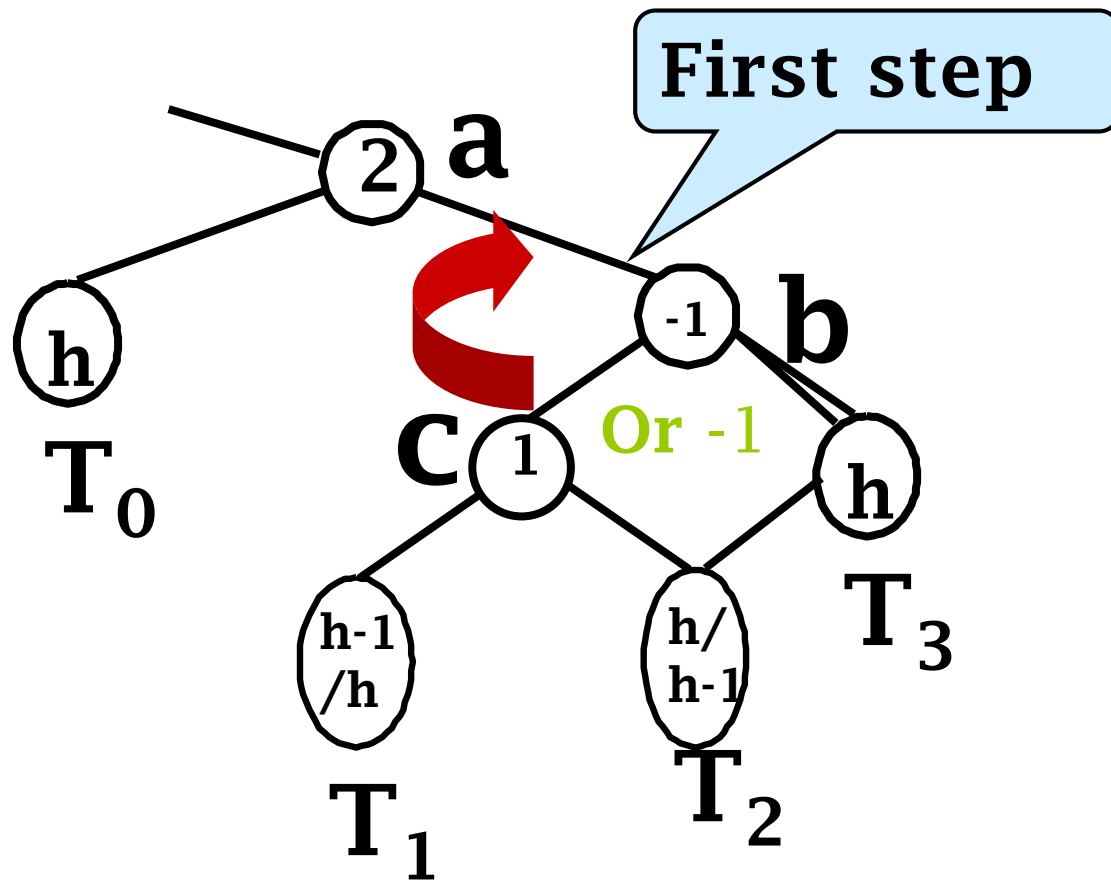
- Take RR for instance, there are 7 parts
 - Three nodes: a, b, c
 - Four subtrees T_0 , T_1 , T_2 , T_3
 - The structure will not change after making c's subtree heavier.
 - T_2 , c, T_3 could be regarded as b's right subtree.
- Goal: construct a new AVL structure
 - Balanced
 - Keep the BST property
 - T_0 a T_1 b T_2 c T_3

Double Rotation

- RL or LR needs double rotation.
 - They are symmetry with each other
- We discuss about RL only
 - LR is the same.

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First step of RL double rotation



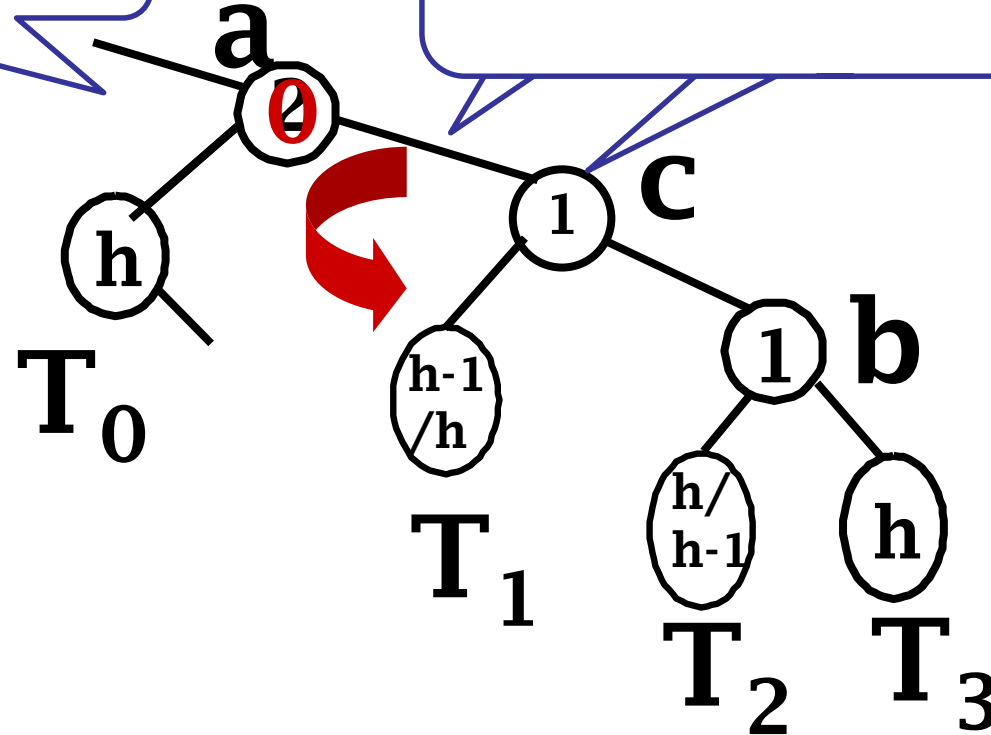
Height of a's subtree is $h+2$ before inserting
 Height of a's subtree is $h+3$ after inserting

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Second step of RL double rotation

Balance factor is meaningless in the middle status

Balance factor of a is -1 or 0
Balance factor of b is 0 or 1



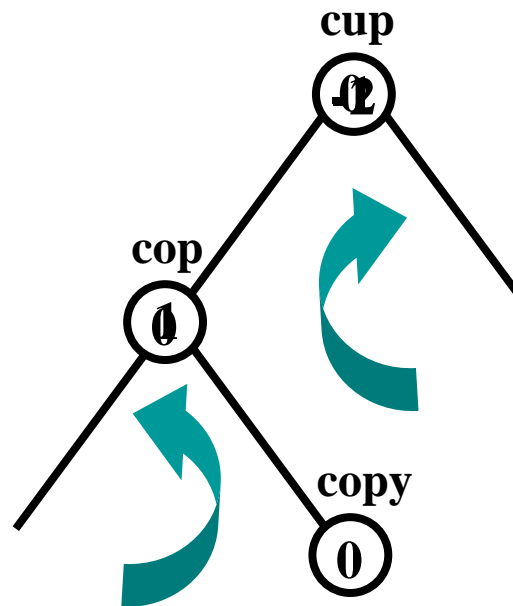


Insight of Rotations

- Doing any rotations (RR, RL, LL, LR)
- New tree keeps the BST property
- Few pointers need to be modified during rotations.
- Height of the new subtree is $h+2$, and heights of subtrees before insertion stay same
- Rest parts upon node a (if not empty) are always balanced

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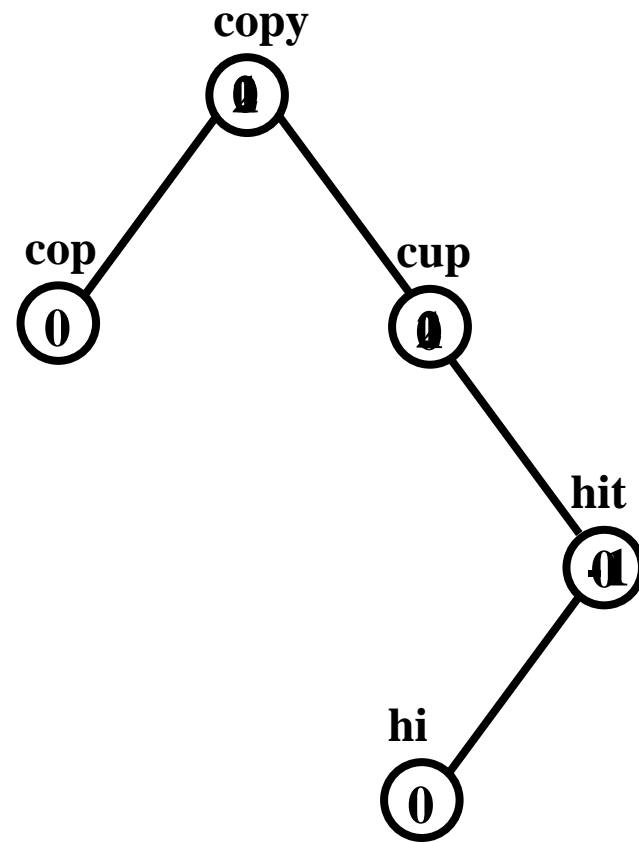
AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia



Unbalanced after
inserting copy
LR double rotation

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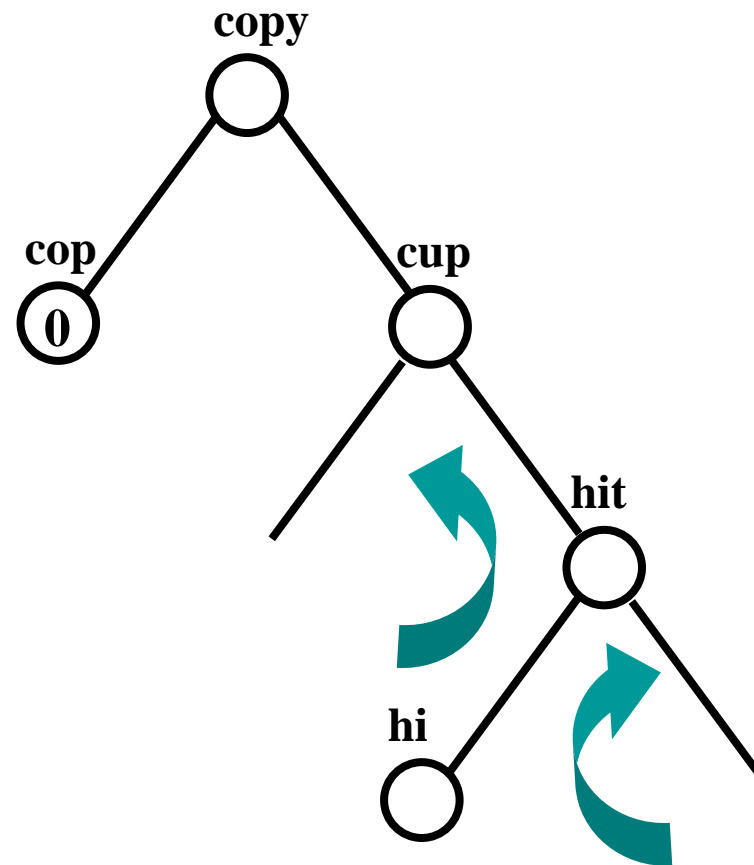
AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia



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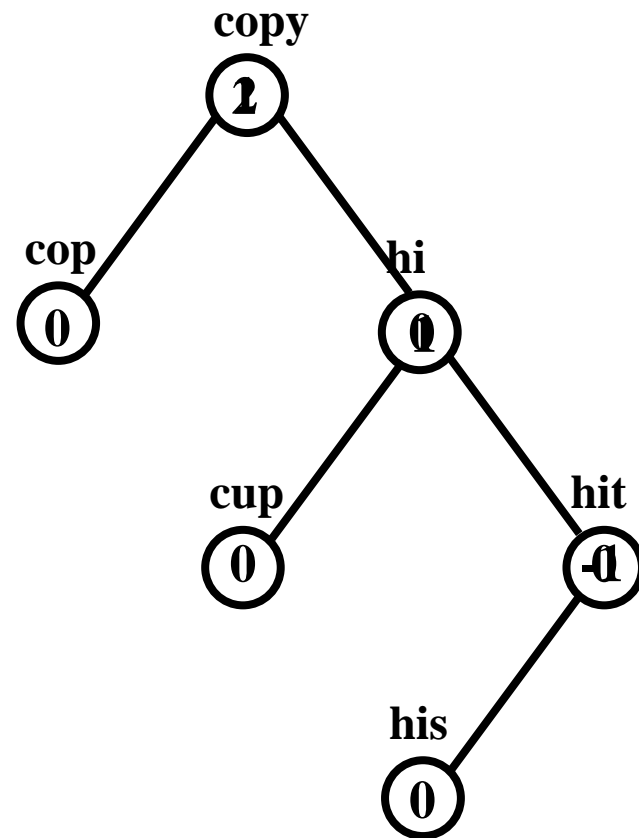
AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia

RL double
rotation



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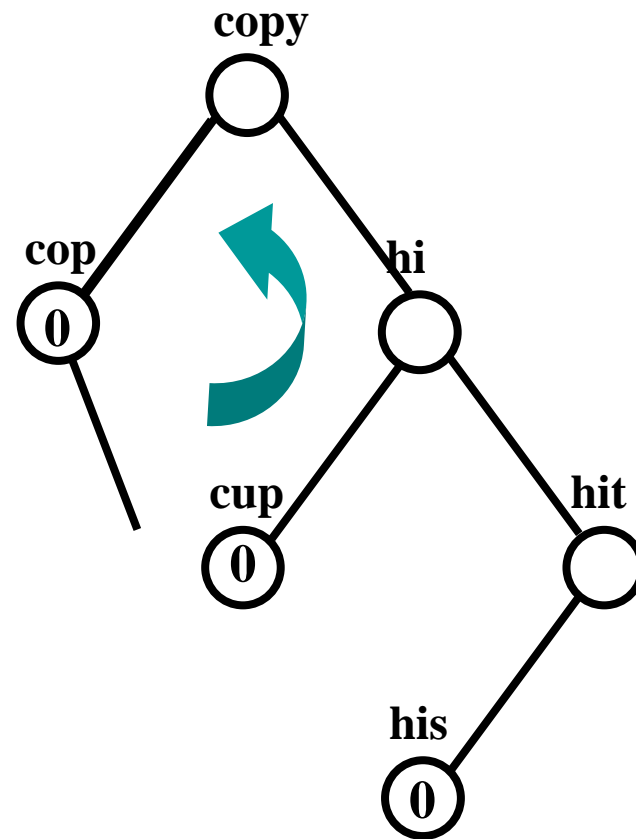
AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia



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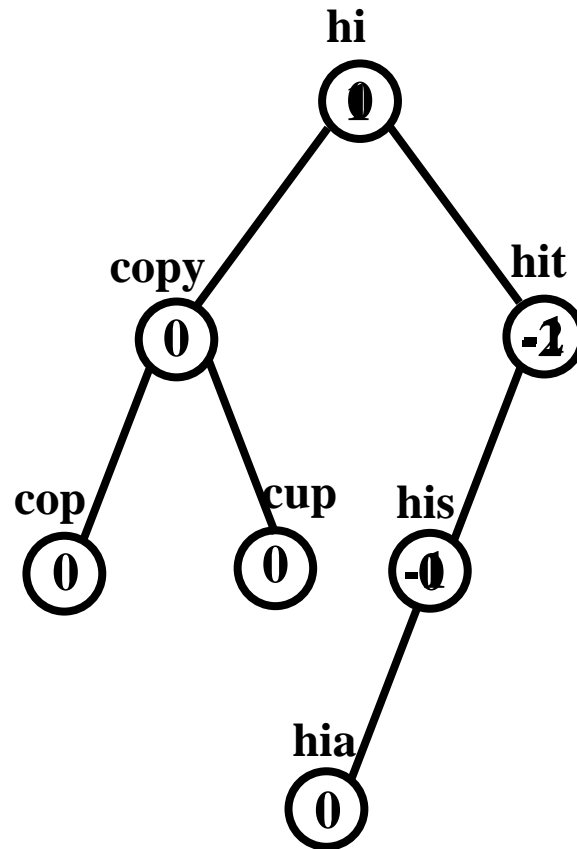
AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia

RR single
rotation



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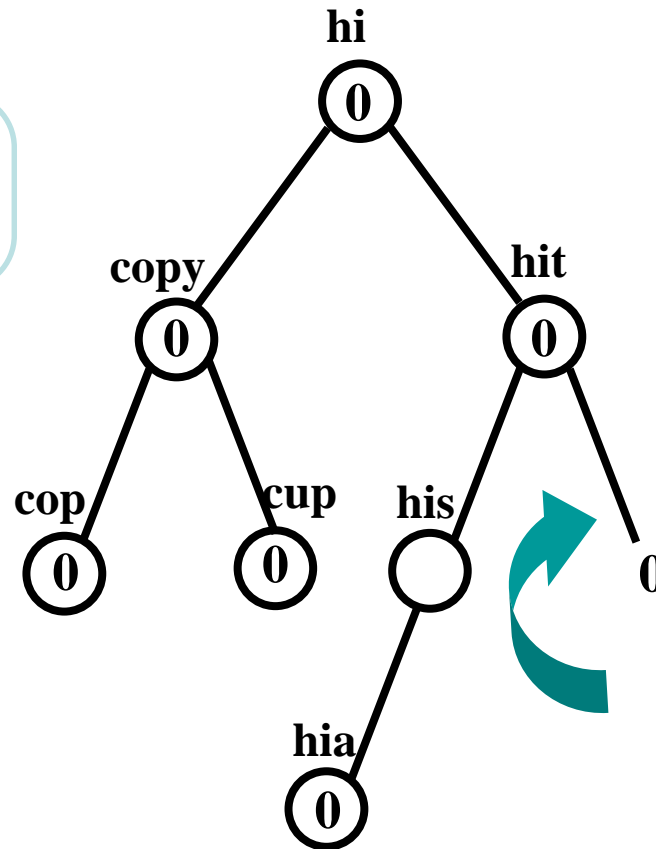
AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia



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AVL tree after inserting word : cup, cop, copy, hit, hi, his and hia

LL single
rotation



12.5 Improved binary search tree

Discussions

- Can we modify the definition of balance factor of AVL tree? For example, allow the height difference as big as 2.
- Insert $1, 2, 3, \dots, 2^k - 1$ into an empty AVL tree consecutively. Try to prove the result is a complete binary tree with height k .



Data Structures and Algorithms

Thanks

the National Elaborate Course (Only available for IPs in China)

<http://www.jpk.pku.edu.cn/pkujpk/course/sjjg/>

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